

## CLAIMS

### What is Claimed is:

1. A method of estimating a communication channel impulse response  $h(t)$ , comprising the steps of:

generating  $co_m(t) = co(t + mNT_c)$  for  $m = 0, 1, \Lambda, M$  by correlating a received signal  $r(t)$  with a spreading sequence  $S_i$  of length  $N$ , wherein the received signal  $r(t)$  comprises a chip sequence  $c_j$  applied to a communication channel characterizable by an impulse response  $h(t)$ , and wherein the chip sequence  $c_j$  is generated from a data sequence  $d_i$  spread by the spreading sequence  $S_i$  and wherein  $T_c$  is the chip period of the chip sequence  $c_j$ ;

generating an estimated communication channel impulse response  $\hat{h}_M(t)$  as a combination of  $co_m(t)$  and  $d_m$  for  $m = 0, 1, \Lambda, M$ ; and

filtering the first estimated communication channel impulse response  $\hat{h}_M(t)$  to generate the estimated communication channel impulse response  $h(t)$  with a filter  $f$  selected at least in part according to the spreading sequence  $S_i$ .

2. The method of claim 1, wherein the filter  $f$  is further selected at least in part according to an autocorrelation  $A(n)$  of the spreading sequence  $S_i$ .

3. The method of claim 2, wherein the filter  $f$  is further selected at least in part according to the duration of the impulse response of the communication channel  $h(t)$ .

4. The method of claim 2, wherein the filter  $f$  is further selected at least in part according to a zero-forcing criteria  $\sum_{i=-L}^L (A(n-i) \bullet f(i)) = A_f(n), -L \leq n \leq L$ , wherein:

$f(i)$  is the impulse response of the filter  $f$  such that  $A_f(n)$  is a convolution of  $A(n)$  and  $f(i)$ ;

$A_f(n) = 1$  for  $n = 0$  and  $A_f(n) = 0$  for  $0 < |n| \leq L$ ; and

$A(n) = A(-n) = \sum_{i=0}^{N-1-n} S_i \bullet S_{i+n}, 0 \leq n \leq N$ , and  $N$  is a length of the chip sequence  $S_i$ .

5. The method of claim 4, wherein:

the parameter  $L$  is chosen such that a time duration of the impulse response of the communication channel  $h(t)$  is less than  $LT_c$ .

6. The method of claim 4, wherein:

the parameter  $L$  is chosen such that a time duration of the impulse response of the communication channel  $h(t)$  is approximately equal to  $LT_c$ .

7. The method of claim 1, wherein  $N$  is less than 20.

8. The method of claim 1, wherein  $M = 0$ .

9. The method of claim 1, wherein the data sequence  $d_i$  includes a constrained portion  $Cd_i$  associated with at least two codes  $w_0, w_1$ , wherein a correlation  $A_{code}(k)$  of the constrained portion  $Cd_i$  with one of the codes  $w_0, w_1$  is characterized by a maximum value at  $k = 0$  less than maximum values at  $k \neq 0$ .

10. The method of claim 9, wherein the step of generating an estimated communication channel impulse response  $\hat{h}_M(t)$  as a combination of  $co_m(t)$  and  $d_m$  for  $m = 0, 1, \Lambda, M$  comprises the step of computing  $\hat{h}_M(t)$  as  $\frac{1}{M} \sum_{m=0}^{M-1} d_m \bullet co(t + mNT_c)$ .

11. The method of claim 10, wherein  $M=2$ .

12. The method of claim 9, wherein the data sequence  $d_i$  includes a preamble having a pseudorandom code including the constrained portion of the data sequence  $d_i$ .

13. The method of claim 9, wherein  $A_{code}(k) = 1$  at  $k = 0$  and  $A_{code}(k) = 0$  for substantially all  $k \neq 0$ .

14. The method of claim 9, wherein  $A_{code}(k) = 0$  for  $0 < |k| \leq J$ , wherein  $J$  is selected to minimize the correlation of the constrained portion  $Cd_i$  with the one of the codes  $w_0, w_1$  for substantially all  $k \neq 0$ .

15. The method of claim 14, wherein  $2J$  is a length of the constrained portion  $Cd_i$ .

16. The method of claim 1, wherein  $A_{code}(k) = 1$  at  $k = 0$  and  $A_{code}(k) = 0$  for substantially all  $k \neq 0$ .

17. The method of claim 1, wherein each of the two codes  $w_0, w_1$  comprises two symbols.

18. The method of claim 1, wherein the each of the two codes  $w_0, w_1$  comprises no more than two symbols.

19. The method of claim 1, wherein the codes  $w_0, w_1$  comprise Walsh codes.

20. An apparatus for estimating a communication channel impulse response  $h(t)$ , comprising:

means for generating  $co_m(t) = co(t + mNT_c)$  for  $m = 0, 1, \Lambda, M$  by correlating a received signal  $r(t)$  with a spreading sequence  $S_i$  of length  $N$ , wherein the received signal  $r(t)$  comprises a chip sequence  $c_j$  applied to a communication channel characterizable by an impulse response  $h(t)$ , and wherein the chip sequence  $c_j$  is generated from a data sequence  $d_i$  spread by the spreading sequence  $S_i$  and wherein  $T_c$  is the chip period of the chip sequence  $c_j$ ;

means for generating an estimated communication channel impulse response  $\hat{h}_M(t)$  as a combination of  $co_m(t)$  and  $d_m$  for  $m = 0, 1, \Lambda, M$ ; and

a filter means  $f$ , selected at least in part according to the spreading sequence  $S_i$ , the filter means for filtering the first estimated communication channel impulse response  $\hat{h}_M(t)$  to generate the estimated communication channel impulse response  $h(t)$  with

21. The apparatus of claim 20, wherein the filter means  $f$  is further selected at least in part according to an autocorrelation  $A(n)$  of the spreading sequence  $S_i$ .

22. The apparatus of claim 21, wherein the filter means  $f$  is further selected at least in part according to the duration of the impulse response of the communication channel  $h(t)$ .

23. The apparatus of claim 21, wherein the filter means  $f$  is further selected at least in part according to a zero-forcing criteria  $\sum_{i=-L}^L (A(n-i) \bullet f(i)) = A_f(n), -L \leq n \leq L$ , wherein:

$f(i)$  is the impulse response of the filter means  $f$  such that  $A_f(n)$  is a convolution of  $A(n)$  and  $f(i)$ ;

$A_f(n) = 1$  for  $n = 0$  and  $A_f(n) = 0$  for  $0 < |n| \leq L$ ; and

$A(n) = A(-n) = \sum_{i=0}^{N-1-n} S_i \bullet S_{i+n}, 0 \leq n \leq N$ , and  $N$  is a length of the chip sequence  $S_i$ .

24. The apparatus of claim 23, wherein:

the parameter  $L$  is chosen such that a time duration of the impulse response of the communication channel  $h(t)$  is less than  $LT_c$ .

25. The apparatus of claim 23, wherein:

the parameter  $L$  is chosen such that a time duration of the impulse response of the communication channel  $h(t)$  is approximately equal to  $LT_c$ .

26. The apparatus of claim 20, wherein  $N$  is less than 20.

27. The apparatus of claim 20, wherein  $M = 0$ .

28. The apparatus of claim 20, wherein the data sequence  $d_i$  includes a constrained portion  $Cd_i$  associated with at least two codes  $w_0, w_1$ , wherein a correlation  $A_{code}(k)$  of the constrained portion  $Cd_i$  with one of the codes  $w_0, w_1$  is characterized by a maximum value at  $k = 0$  less than maximum values at  $k \neq 0$ .

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29. The apparatus of claim 28, wherein the means for generating an estimated communication channel impulse response  $\hat{h}_M(t)$  as a combination of  $co_m(t)$  and  $d_m$  for  $m = 0, 1, \Lambda, M$  comprises means for computing  $\hat{h}_M(t)$  as  $\frac{1}{M} \sum_{m=0}^{M-1} d_m \bullet co(t + mNT_c)$ .

30. The apparatus of claim 29, wherein  $M=2$ .

31. The apparatus of claim 28, wherein the data sequence  $d_i$  includes a preamble having a pseudorandom code including the constrained portion of the data sequence  $d_i$ .

32. The apparatus of claim 28, wherein  $A_{code}(k) = 1$  at  $k = 0$  and  $A_{code}(k) = 0$  for substantially all  $k \neq 0$ .

33. The apparatus of claim 28, wherein  $A_{code}(k) = 0$  for  $0 < |k| \leq J$ , wherein  $J$  is selected to minimize the correlation of the constrained portion  $Cd_i$  with the one of the codes  $w_0, w_1$  for substantially all  $k \neq 0$ .

34. The apparatus of claim 33, wherein  $2J$  is a length of the constrained portion  $Cd_i$ .

35. The apparatus of claim 20, wherein  $A_{code}(k) = 1$  at  $k = 0$  and  $A_{code}(k) = 0$  for substantially all  $k \neq 0$ .

36. The apparatus of claim 20, wherein each of the two codes  $w_0, w_1$  comprises two symbols.

37. The apparatus of claim 20, wherein the each of the two codes  $w_0, w_1$  comprises no more than two symbols.

38. The apparatus of claim 20, wherein the codes  $w_0, w_1$  comprise Walsh codes.

39. An apparatus for estimating a communication channel impulse response  $h(t)$ , comprising:

a correlator generating  $co_m(t) = co(t + mNT_c)$  for  $m = 0, 1, \Lambda, M$  by correlating a received signal  $r(t)$  with a spreading sequence  $S_i$  of length  $N$ , wherein the received signal  $r(t)$  comprises a chip sequence  $c_j$  applied to a communication channel characterizable by an impulse response  $h(t)$ , and wherein the chip sequence  $c_j$  is generated from a data sequence  $d_i$  spread by the spreading sequence  $S_i$  and wherein  $T_c$  is the chip period of the chip sequence  $c_j$ ;

an estimator for generating an estimated communication channel impulse response  $\hat{h}_M(t)$  as a combination of  $co_m(t)$  and  $d_m$  for  $m = 0, 1, \Lambda, M$ ; and

a filter  $f$  selected at least in part according to the spreading sequence  $S_i$ , the filter for filtering the first estimated communication channel impulse response  $\hat{h}_M(t)$  to generate the estimated communication channel impulse response  $h(t)$ .

40. The apparatus of claim 39, wherein the filter  $f$  is further selected at least in part according to an autocorrelation  $A(n)$  of the spreading sequence  $S_i$ .

41. The apparatus of claim 40, wherein the filter  $f$  is further selected at least in part according to the duration of the impulse response of the communication channel  $h(t)$ .

42. The apparatus of claim 40, wherein the filter  $f$  is further selected at least in part according to a zero-forcing criteria  $\sum_{i=-L}^L (A(n-i) \bullet f(i)) = A_f(n), -L \leq n \leq L$ , wherein:

$f(i)$  is the impulse response of the filter  $f$  such that  $A_f(n)$  is a convolution of  $A(n)$  and  $f(i)$ ;

$A_f(n) = 1$  for  $n = 0$  and  $A_f(n) = 0$  for  $0 < |n| \leq L$ ; and

$A(n) = A(-n) = \sum_{i=0}^{N-1-n} S_i \bullet S_{i+n}, 0 \leq n \leq N$ , and  $N$  is a length of the chip sequence  $S_i$ .

43. The apparatus of claim 42, wherein:

the parameter  $L$  is chosen such that a time duration of the impulse response of the communication channel  $h(t)$  is less than  $LT_c$ .

44. The apparatus of claim 42, wherein:

the parameter  $L$  is chosen such that a time duration of the impulse response of the communication channel  $h(t)$  is approximately equal to  $LT_c$ .

45. The apparatus of claim 39, wherein  $N$  is less than 20.

46. The apparatus of claim 39, wherein  $M = 0$ .

47. The apparatus of claim 39, wherein the data sequence  $d_i$  includes a constrained portion  $Cd_i$  associated with at least two codes  $w_0, w_1$ , wherein a correlation  $A_{code}(k)$  of the constrained portion  $Cd_i$  with one of the codes  $w_0, w_1$  is characterized by a maximum value at  $k = 0$  less than maximum values at  $k \neq 0$ .



48. The apparatus of claim 47, wherein the estimator for generating an estimated communication channel impulse response  $\hat{h}_M(t)$  as a combination of  $co_m(t)$  and  $d_m$  for  $m = 0, 1, \Lambda, M$  comprises means for computing  $\hat{h}_M(t)$  as

$$\frac{1}{M} \sum_{m=0}^{M-1} d_m \bullet co(t + mNT_c).$$

49. The apparatus of claim 48, wherein  $M=2$ .

50. The apparatus of claim 47, wherein the data sequence  $d_i$  includes a preamble having a pseudorandom code including the constrained portion of the data sequence  $d_i$ .

51. The apparatus of claim 47, wherein  $A_{code}(k) = 1$  at  $k = 0$  and  $A_{code}(k) = 0$  for substantially all  $k \neq 0$ .

52. The apparatus of claim 47, wherein  $A_{code}(k) = 0$  for  $0 < |k| \leq J$ , wherein  $J$  is selected to minimize the correlation of the constrained portion  $Cd_i$  with the one of the codes  $w_0, w_1$  for substantially all  $k \neq 0$ .

53. The apparatus of claim 52, wherein  $2J$  is a length of the constrained portion  $Cd_i$ .

54. The apparatus of claim 39, wherein  $A_{code}(k) = 1$  at  $k = 0$  and  $A_{code}(k) = 0$  for substantially all  $k \neq 0$ .

55. The apparatus of claim 39, wherein each of the two codes  $w_0, w_1$  comprises two symbols.

56. The apparatus of claim 39, wherein the each of the two codes  $w_0, w_1$  comprises no more than two symbols.

57. The apparatus of claim 39, wherein the codes  $w_0, w_1$  comprise Walsh codes.